

1.4 Systems of Linear Equations

Question: When is \vec{b} found in $\text{Span}(\{\vec{v}_1, \vec{v}_2, \dots, \vec{v}_n\})$?

Example: Let $\vec{b} = \langle -6, 4, 8 \rangle$. Is \vec{b} in:

- L : The line $\text{Span}(\{\langle 9, -6, -12 \rangle\})$?
- Π_1 : The plane with Cartesian equation $2x - 4y + 5z = 0$?
- Π_2 : The plane with Cartesian equation $2x + 9y - 3z = 0$?

Example: Decide if $\vec{b} = \langle 10, -9, -5, -7 \rangle$ is a member of the Span of the following five vectors from \mathbb{R}^4 :

$$\langle 3, -4, 1, -6 \rangle, \langle 2, -3, 2, -5 \rangle, \langle 1, 1, -9, 5 \rangle,$$
$$\langle 1, -2, 2, -4 \rangle, \langle 9, -7, -8, -3 \rangle$$

If so, express \vec{b} as a linear combination of these five vectors in the simplest way possible.

Follow up:

Matrices

Definition: A *matrix* is a rectangular table of numbers organized into m *rows* and n *columns*. We say that the *dimension* of the matrix is $m \times n$, pronounced “ m *by* n .” If A is an $m \times n$ matrix, we call the number or *entry* in row i column j as $(A)_{i,j}$ or $a_{i,j}$.

Thus we write:

$$A = \begin{bmatrix} a_{1,1} & a_{1,2} & \cdots & a_{1,n} \\ a_{2,1} & a_{2,2} & \cdots & a_{2,n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m,1} & a_{m,2} & \cdots & a_{m,n} \end{bmatrix}.$$

In particular, an $m \times 1$ matrix will be called a *column matrix*, and a $1 \times n$ matrix will be called a *row matrix*.

If $S = \{ \vec{v}_1, \vec{v}_2, \dots, \vec{v}_n \}$ is a set of vectors from \mathbb{R}^m , and $\vec{b} \in \mathbb{R}^m$, we can form the $m \times (n + 1)$ *augmented matrix*:

$$A = \left[\vec{v}_1 \quad \vec{v}_2 \quad \dots \quad \vec{v}_n \quad | \quad \vec{b} \right],$$

where we assemble the vectors in S into *columns*, and we separate the last column \vec{b} with a dashed line to indicate that it represents the right side of a system of equations. We normally use x_1, x_2, \dots, x_n to be the variables *associated* to each column.

The Reduced Row Echelon Form

or RREF

Definition: We will say that an $m \times (n + 1)$ augmented matrix is in *row echelon form* if it satisfies the following conditions:

1. All the rows consisting entirely of zeroes are at the *bottom* of the matrix.
2. The first non-zero entry of any row is the number 1. This entry is called a “*leading 1*.”
3. If the next row is non-zero, its leading 1 is to the *right* of the previous leading 1.

Furthermore, we say that the matrix is in *reduced row echelon form*, or *rref*, if the matrix satisfies the additional condition:

4. All the entries *above* a leading 1 are also zeroes.

If column j contains a leading 1, we call x_j a *leading variable*, and column j is a *leading column*.

Otherwise, we call x_j a *free variable*.

Conditions 1 and 3 forces all entries *below* a leading 1 to be zeroes. Condition 4 forces all entries *above* the leading 1 to be zeroes as well.

The vectors $\vec{e}_1, \vec{e}_2, \dots, \vec{e}_k$ will appear in the leading columns, in that order, for some $k \leq m$.

Example:

$$A = \left[\begin{array}{cccccc|c} 1 & 0 & -4 & 0 & 9 & 2 & 6 \\ 0 & 1 & 3 & 0 & 6 & -5 & -9 \\ 0 & 0 & 0 & 1 & -2 & 7 & -4 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array} \right]$$

Elementary Row Operations

Definition: An *elementary row operation* is any one of the following actions on a matrix:

Type:

Notation:

1. Multiply row i by a nonzero scalar c

$$R_i \rightarrow cR_i.$$

2. Exchange row i and row j

$$R_i \leftrightarrow R_j.$$

3. Add c times row j to row i

$$R_i \rightarrow R_i + cR_j.$$

Invariance of Solution Sets

Theorem — The Invariance of Solution Sets:

An elementary row operation *does not change* the solution set of an augmented matrix. In other words, if A is an augmented matrix and B is obtained from A using an elementary row operation, then the solution set of the system corresponding to A is exactly the same as the solution set of the system corresponding to B .

The Uniqueness of the RREF

Theorem — The Uniqueness of the Reduced Row Echelon Form:

The reduced row echelon form of a matrix is *unique*. Thus, if we start with a matrix A and arrive at two matrices H and J using a different sequence of row operations, and both H and J are in rref, then $H = J$.

The Gauss-Jordan Algorithm

The Gauss-Jordan Algorithm

1. Ignore all the leftmost columns that contain only zeros, if there are any.
2. Starting from the top row and going downward, find the first non-zero entry, called the *pivot*.
3. If the pivot is not in the top row, *exchange* the top row with the pivot's row (this is a Type 2 row operation).
4. Produce a leading 1 in the top row by *dividing* the entire top row by the pivot (this is a Type 1 row operation). We call this step *normalizing* the row.
5. Make the entries below the leading 1 *all zeroes* by adding suitable multiples of the top row to each row below it (these are Type 3 row operations).
6. Now, cover the top row, the leading column and all columns to its left, and repeat steps 1 through 5 on the smaller submatrix.

If we were to stop at Step 6, the algorithm above would be called *Gaussian Elimination*. It results in a matrix in *row echelon form*. Now we continue from *right to left*, working *upwards* as we go:

7. Starting at the rightmost leading 1, produce zeroes above the leading 1 by adding suitable multiples of this row to each row above it (again, these are Type 3 row operations).

8. Repeat Step 7 on the next rightmost leading 1, moving leftward until the matrix is in *reduced row echelon form*.

An Intelligent Modification:

To avoid fractions, produce a leading “1” using a Type 3 operation instead: multiply a row by a suitable constant and add it to another row to get a leading 1.

Success and Failure

\vec{b} is *not* in $\text{Span}(S)$ if we get a row consisting entirely of zeroes except for a non-zero entry in the rightmost column.

$$\left[\begin{array}{ccccc|c} 1 & -7 & 0 & 6 & 0 & 3 \\ 0 & 0 & 1 & -4 & 0 & 2 \\ 0 & 0 & 0 & 0 & 1 & -5 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{array} \right]$$